

# Decentralized Finance

## Introduction to Blockchain technology

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# What is a blockchain?

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Abstract answer: a blockchain provides coordination between many parties, when there is no single trusted party

if trusted party exists  $\Rightarrow$  no need for a blockchain

[financial systems: often no trusted party]

# What is a blockchain?

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**user facing tools** (cloud servers)

**applications** (DAPPs, smart contracts)

**compute layer** (blockchain computer)

**consensus layer**

# Consensus layer (informal – not the topic of this course)

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A public append-only data structure:

achieved by replication

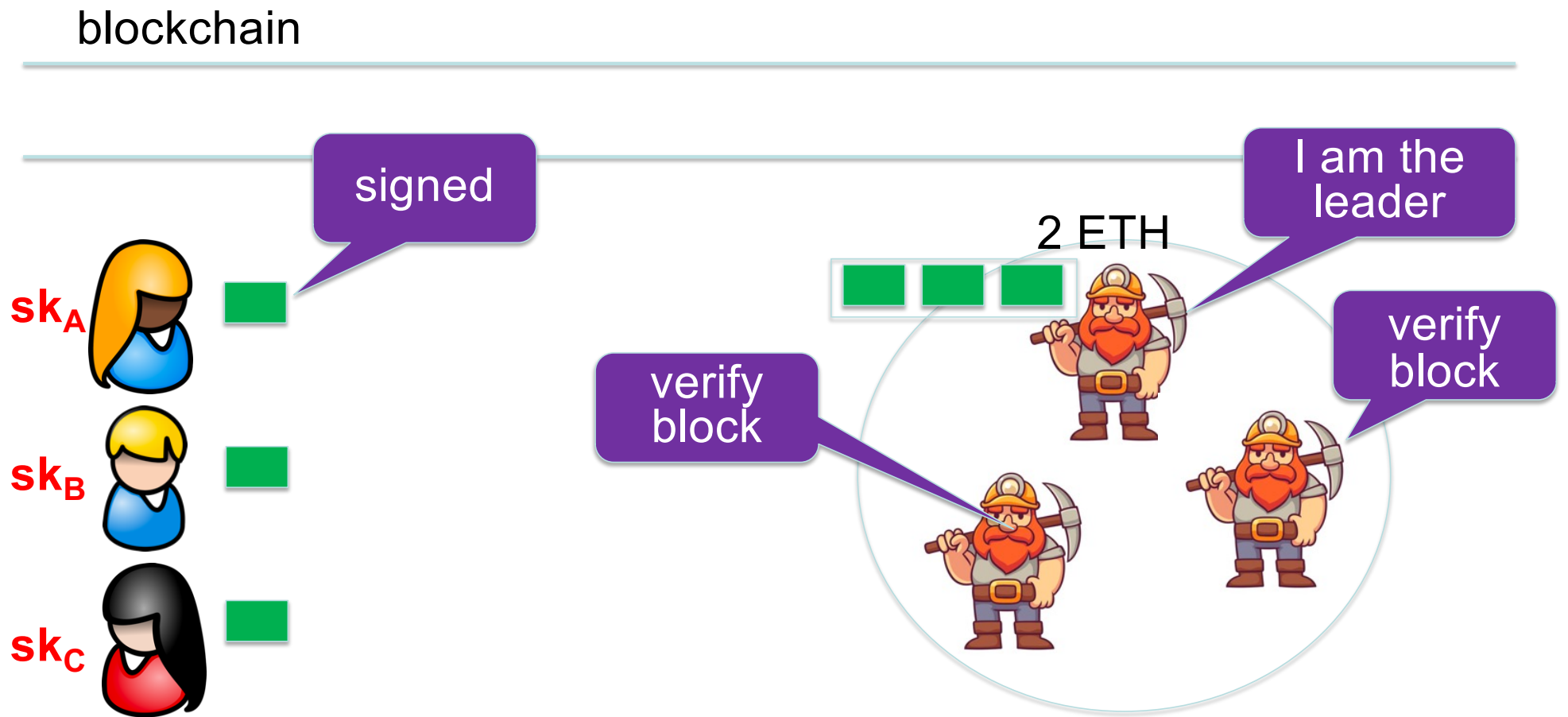


- **Persistence:** once added, data can never be removed\*
- **Consensus:** all honest participants have the same data\*\*
- **Liveness:** honest participants can add new transactions
- **Open(?):** anyone can add data

Layer 1:

consensus layer

# How are blocks added to chain?

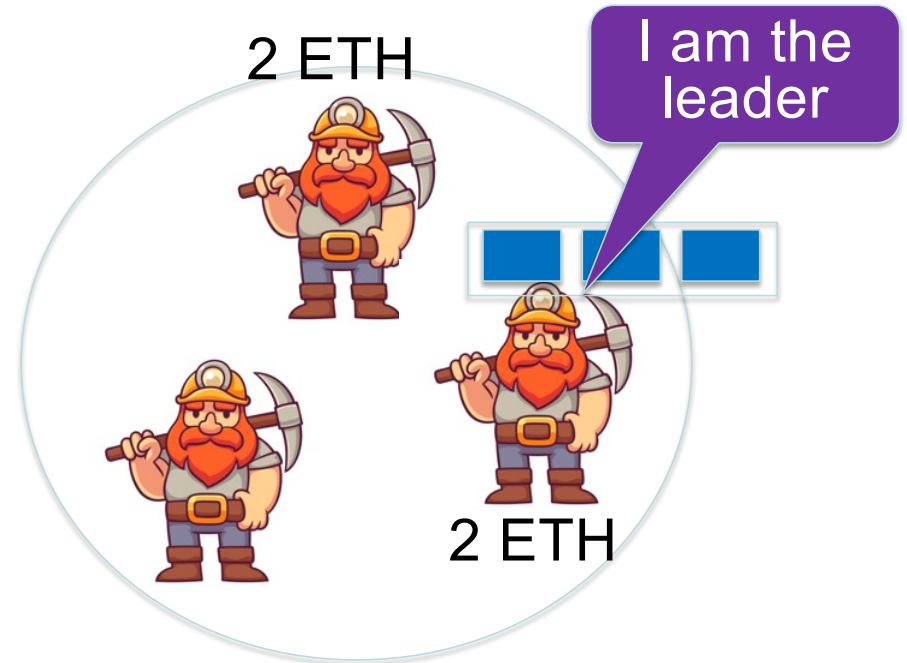
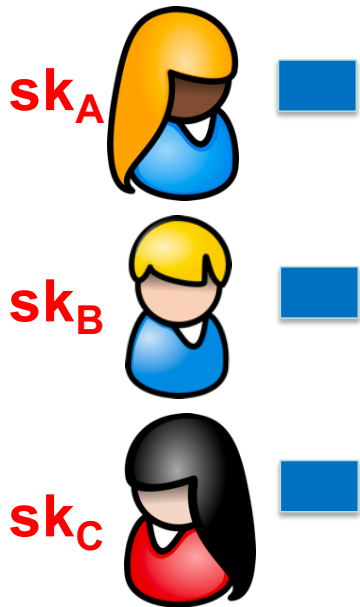


# How are blocks added to chain?

blockchain



...



# Compute layer: The blockchain computer

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## DAPP logic is encoded in a program that runs on blockchain

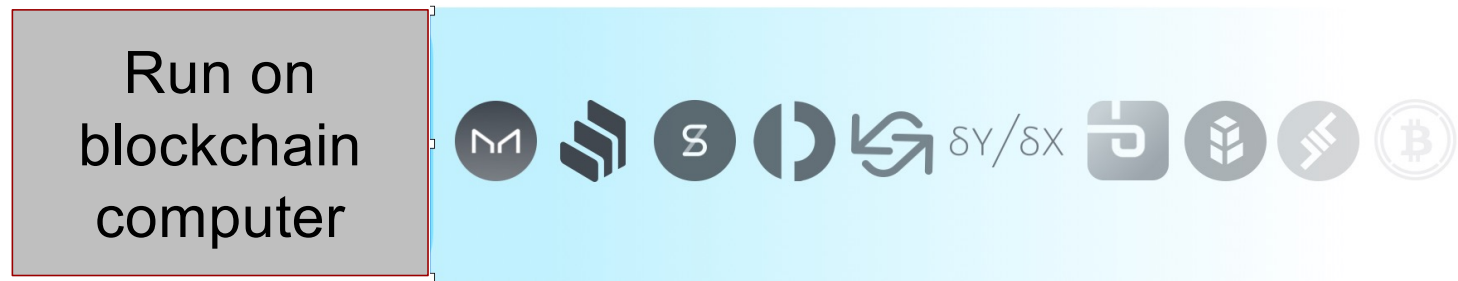
- Rules are enforced by a public program (public source code)
  - ⇒ transparency: no single trusted 3<sup>rd</sup> party
- The DAPP program is executed by parties who create new blocks
  - ⇒ public verifiability: everyone can verify state transitions

compute layer

consensus layer

# Apps layer: Decentralized applications (DAPPS)

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**applications** (DAPPs, smart contracts)

blockchain computer

consensus layer

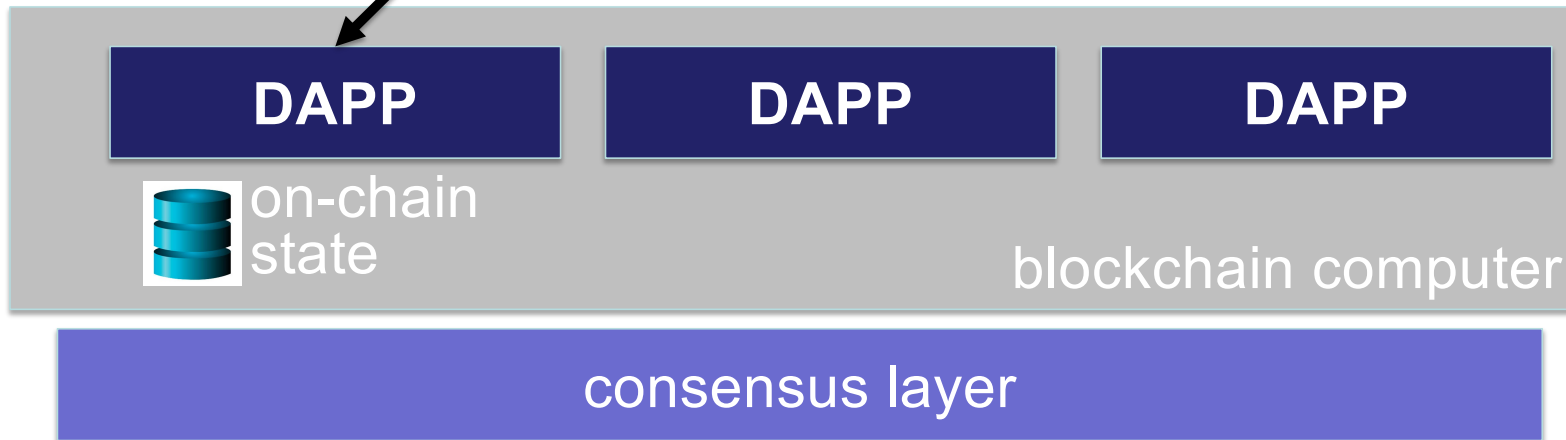


# UI Layer: Common DAPP architecture

UI Layer: user facing servers



end user



# Ethereum's DeFi



## Payments



## OPEN PLATFORM



## Custodial Services



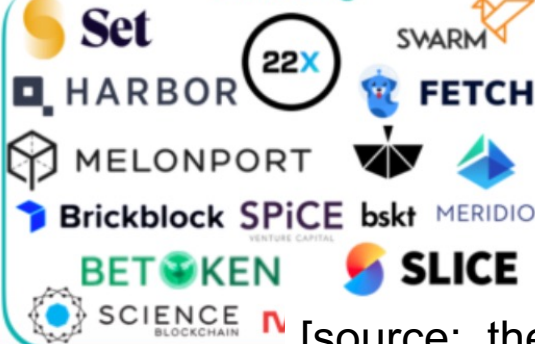
## Infrastructure



## Exchanges & Liquidity



## Investing



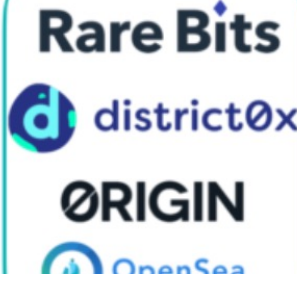
## KYC & Identity



## Derivatives



## Marketplaces



## Stablecoins



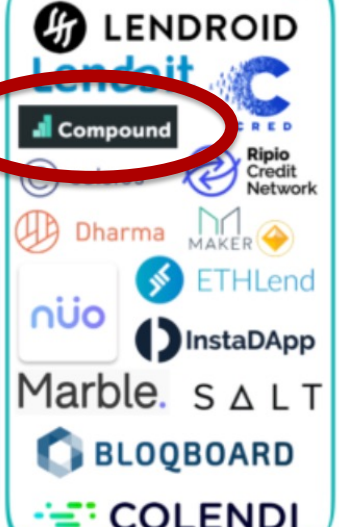
## Prediction Markets



## Insurance






## Credit & Lending



[source: the Block Genesis]

# lots of experiments ...

DEFI PULSE	Name	Chain	Category	Locked (USD) ▼
 1.	Aave	Multichain	Lending	\$15.63B
 2.	Curve Finance	Multichain	DEXes	\$10.71B
 3.	InstaDApp	Ethereum	Lending	\$10.66B
4.	Compound	Ethereum	Lending	\$10.43B
5.	Maker	Ethereum	Lending	\$9.11B
6.	Uniswap	Ethereum	DEXes	\$7.29B
7.	Convex Finance	Ethereum	Assets	\$5.59B

Let's get started ...

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Next segment: cryptographic background

See you there



# Cryptographic Background: hash functions

<https://defi-learning.org/>

# Cryptography Background

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## (1) cryptographic hash functions

An efficiently computable function  $H: M \rightarrow T$

where  $|M| \gg |T|$



# Collision resistance

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**Def:** a collision for  $H: M \rightarrow T$  is pair  $x \neq y \in M$  s.t.  $H(x) = H(y)$

$|M| \gg |T|$  implies that many collisions exist

**Def:** a function  $H: M \rightarrow T$  is collision resistant if it is “hard” to find even a single collision for  $H$  (we say  $H$  is a CRH)

Example: **SHA256:**  $\{x : \text{len}(x) < 2^{64} \text{ bytes}\} \rightarrow \{0,1\}^{256}$

details in crypto MOOC

# An application: committing to data

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Alice has a large file  $m$ . She publishes  $h = H(m)$  (32 bytes)

Bob has  $h$ . Later Alice sends  $m'$  s.t.  $H(m') = h$

$H$  is a CRH  $\Rightarrow$  Bob is convinced that  $m' = m$   
(otherwise,  $m$  and  $m'$  are a collision for  $H$ )

We say that  $h = H(m)$  is a **binding commitment** to  $m$

(note: not hiding,  $h$  may leak information about  $m$ )



# Committing to a list (of transactions)

Alice has  $S = (m_1, m_2, \dots, m_n)$

32 bytes



## Goal:

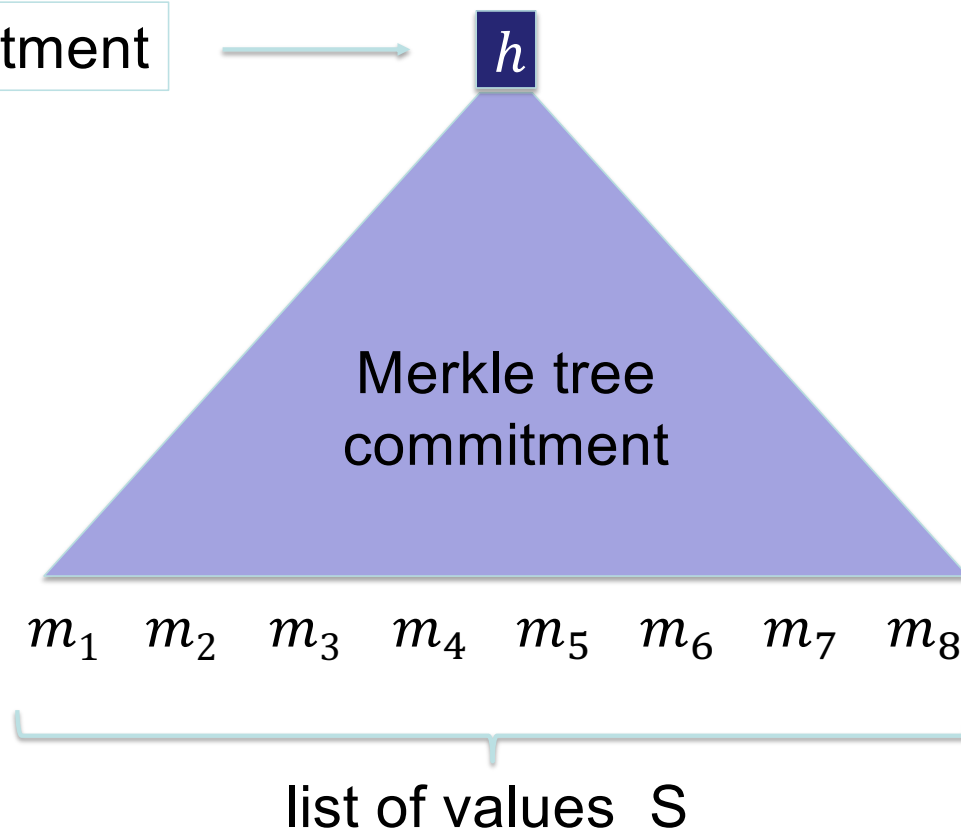
- Alice publishes a short binding commitment to  $S$ ,  $h = \text{commit}(S)$
- Bob has  $h$ . Given  $(m_i, \text{proof } \pi_i)$  can check that  $S[i] = m_i$   
Bob runs  $\text{verify}(h, i, m_i, \pi_i) \rightarrow \text{accept/reject}$

security: adv. cannot find  $(S, i, m, \pi)$  s.t.  $m \neq S[i]$  and

$\text{verify}(h, i, m, \pi) = \text{accept}$  where  $h = \text{commit}(S)$

# Merkle tree (Merkle 1989)

commitment

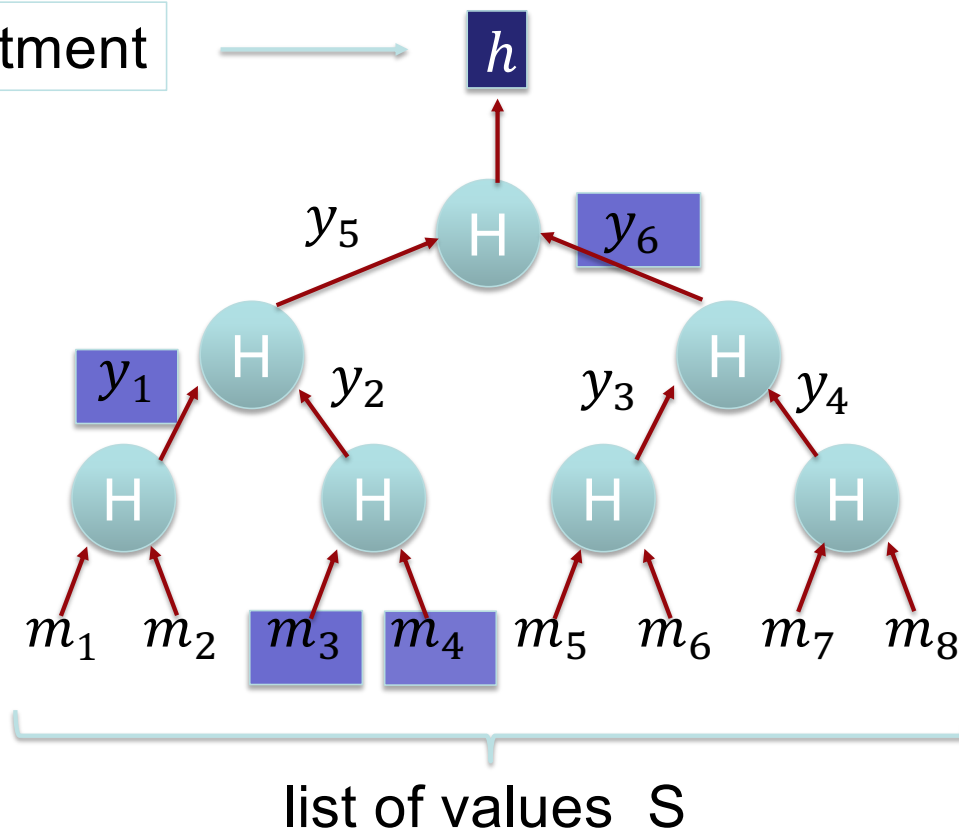


Goal:

- commit to list  $S$  of size  $n$
- Later prove  $S[i] = m_i$

# Merkle tree (Merkle 1989)

commitment



Goal:

- commit to list  $S$  of size  $n$
- Later prove  $S[i] = m_i$

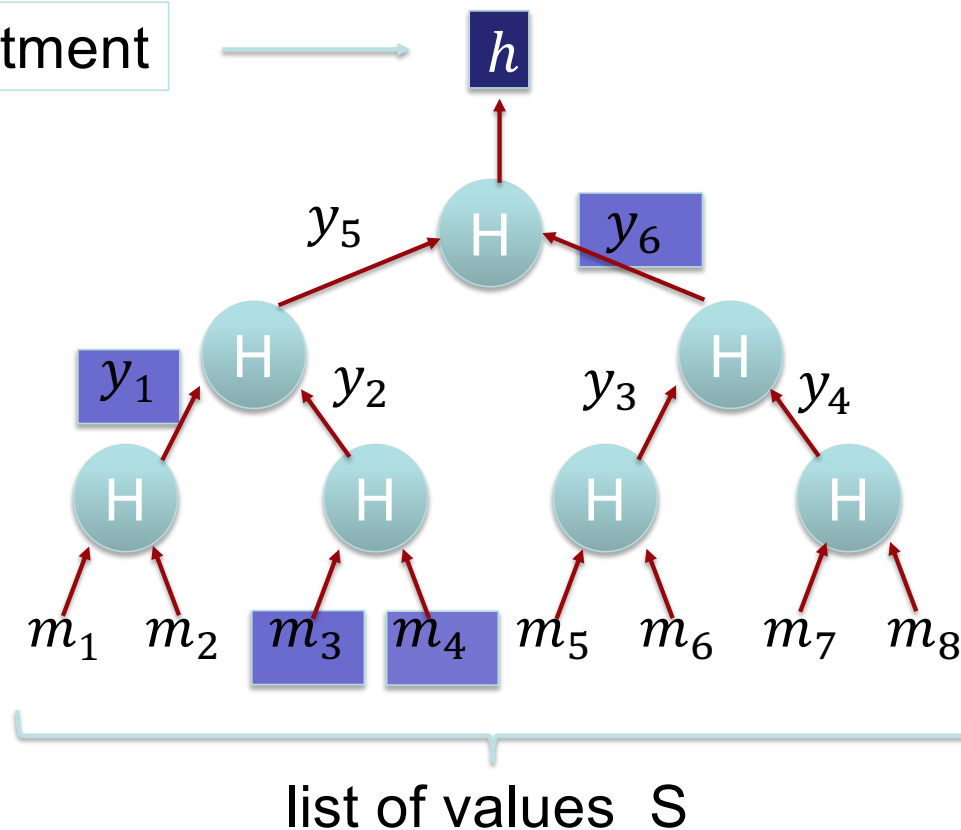
To prove  $S[4] = m_4$

proof  $\pi = (m_3, y_1, y_6)$

length of  $\pi$ :  $\log_2 n$

# Merkle tree (Merkle 1989)

commitment



To prove  $S[4] = m_4$

proof  $\pi = (m_3, y_1, y_6)$

Bob does:

$$y_2 \leftarrow H(m_3, m_4)$$

$$y_5 \leftarrow H(y_1, y_2)$$

$$h' \leftarrow H(y_5, y_6)$$

accept if  $h = h'$

# Merkle tree (Merkle 1989)

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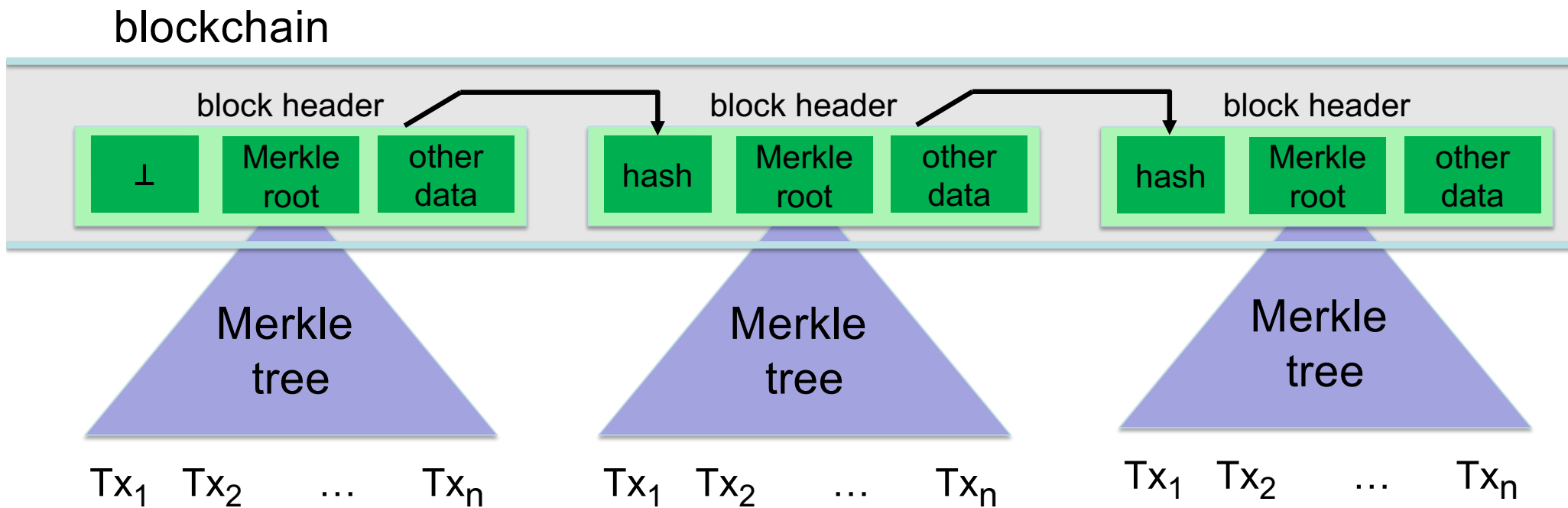
Thm:  $H$  CRH  $\Rightarrow$  adv. cannot find  $(S, i, m, \pi)$  s.t.  $m \neq S[i]$  and  
 $\text{verify}(h, i, m, \pi) = \text{accept}$  where  $h = \text{commit}(S)$

(to prove, prove the contra-positive)

How is this useful? Super useful. Example:

- When writing a block of transactions  $S$  to the blockchain, suffices to write  $\text{commit}(S)$  to chain. Keep chain small.
- Later, can prove contents of every Tx.

# Abstract block chain



Merkle proofs are used to prove that a Tx is “on the block chain”

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Next segment: digital signatures

How to authorize transactions??



# Cryptographic Background: Digital Signatures

<https://defi-learning.org/>



# Digital Signatures

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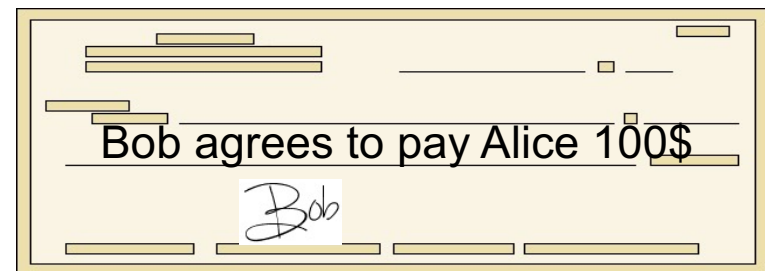
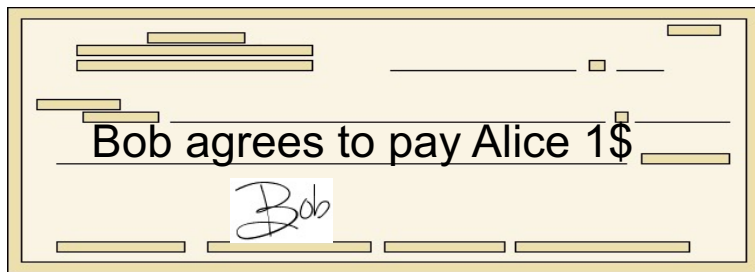
- In the last segment we looked at cryptographic hash functions.
- In this segment we will look at digital signatures:

how to approve a transaction?

# Signatures

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Physical signatures: bind transaction to author

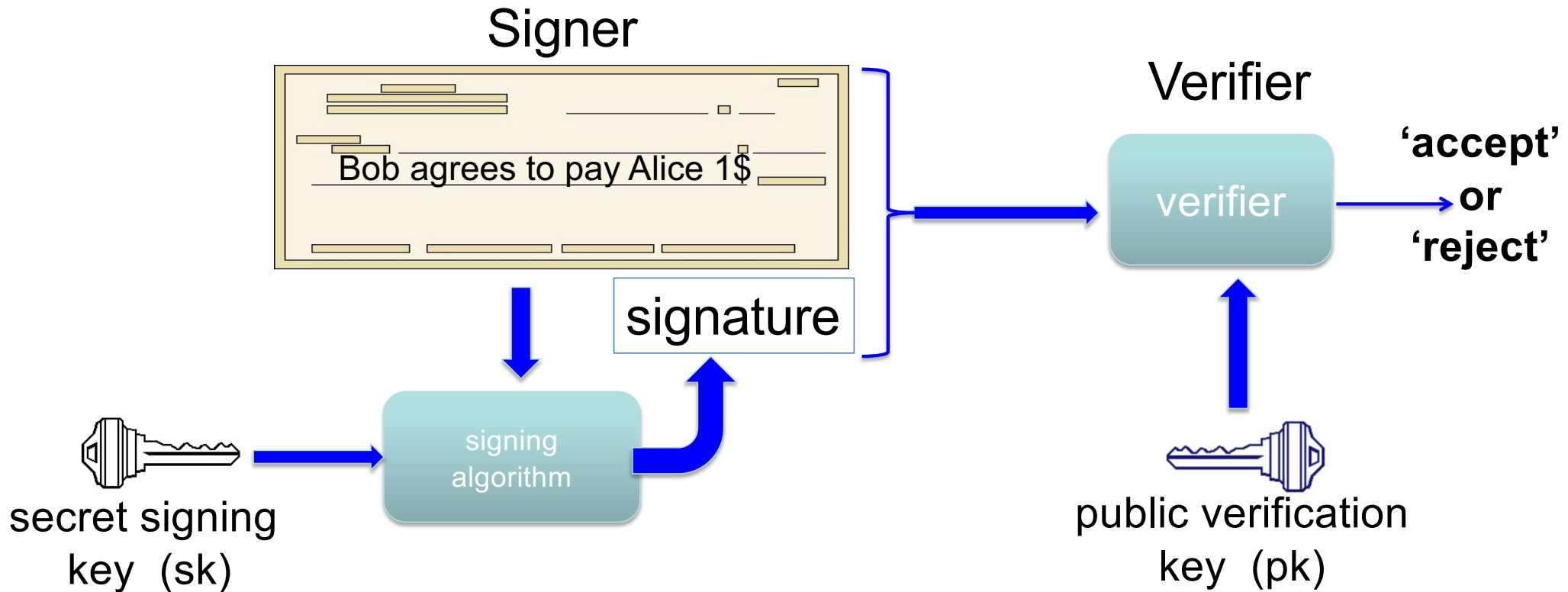


Problem in the digital world:

anyone can copy Bob's signature from one doc to another

# Digital signatures

Solution: make signature depend on document



# Digital signatures: syntax

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Def: a signature scheme is a triple of algorithms:

- **Gen()**: outputs a key pair  $(pk, sk)$
- **Sign**(sk, msg) outputs sig.  $\sigma$
- **Verify**(pk, msg,  $\sigma$ ) outputs 'accept' or 'reject'

Secure signatures: (informal)

Adversary who sees pk and sigs on many messages of her choice, cannot forge a signature on a new message.

# Families of signature schemes

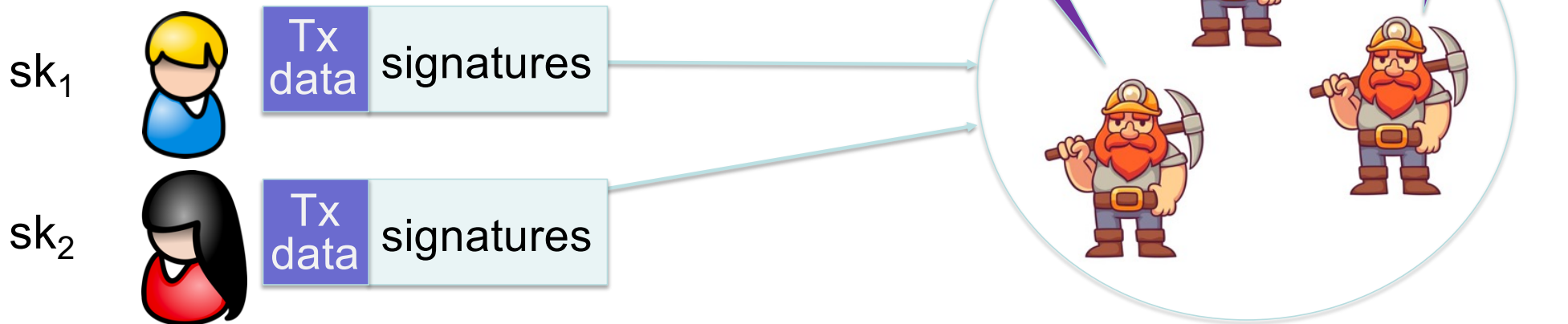
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1. RSA signatures (not used in blockchains):
    - long sigs and public keys ( $\geq 256$  bytes), fast to verify
  2. Discrete-log signatures: Schnorr and ECDSA (Bitcoin, Ethereum)
    - short sigs (48 or 64 bytes) and public keys (32 bytes)
  3. BLS signatures: 48 bytes, aggregatable, easy threshold  
(Ethereum 2.0, Chia, Dfinity)
- 
4. Post-quantum signatures: long ( $\geq 768$  bytes)

# Signatures on the blockchain

Signatures are used everywhere:

- ensure Tx authorization,
- governance votes,
- consensus protocol votes.



# SNARK proofs

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We covered two important cryptographic primitives:

1. Collision resistant hash functions and Merkle trees,
2. Digital signatures.

Another important cryptographic primitive is a **SNARK proof**:

- Used for scaling and privacy
- We will discuss SNARKs in detail in the lecture on privacy

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Next segment: scaling the blockchains

Can we make it fast??





# Scaling Blockchains

<https://defi-learning.org/>

# Scaling

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Transaction rates (Tx/sec):

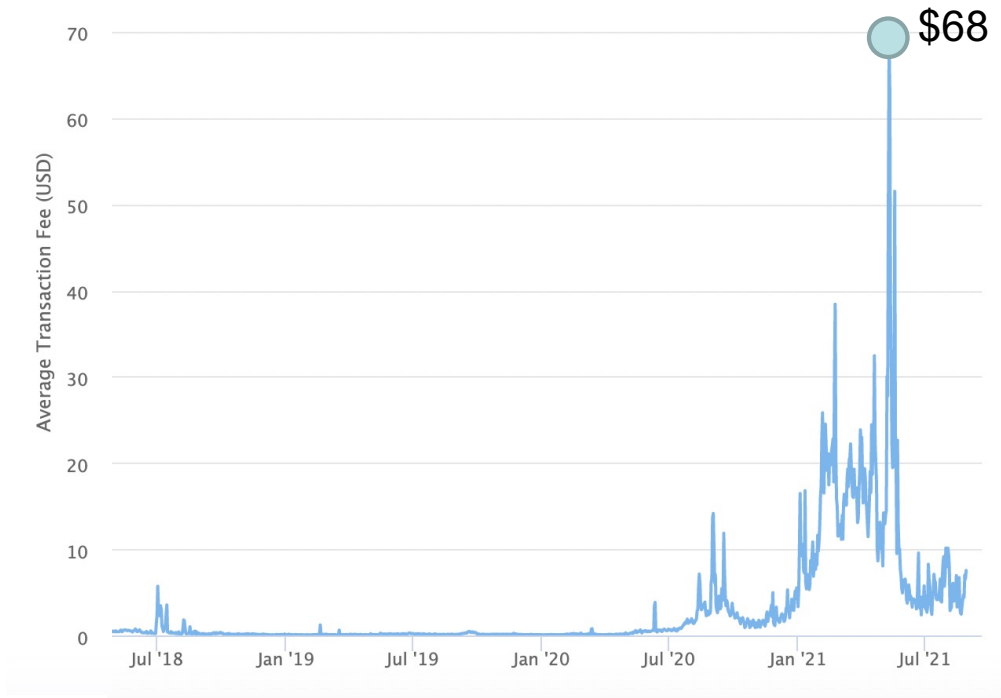
- Bitcoin: can process about **5 (Tx/sec)**
- Ethereum: can process about **20 (Tx/sec)**

} Tx Fees fluctuate:  
2\$ to 60\$ for simple Tx

# Ethereum Tx fees (gas prices)

Average Transaction Fee Chart

Source: Etherscan.io



# Scaling

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Transaction rates (Tx/sec):

- Bitcoin: can process about **5 (Tx/sec)**
  - Ethereum: can process about **20 (Tx/sec)**
  - The visa network: can process up to **24,000 (Tx/sec)**
- Tx Fees fluctuate:  
2\$ to 60\$ for simple Tx

**Can we scale blockchains to visa speeds? ... with low Tx fees**

# Scaling approaches

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Many approaches to scaling blockchains:

- **Faster consensus:** modern blockchains (e.g., Solana, Polkadot, Avalanche, ...)
- **Payment channels:** most Tx are off chain Peer-to-Peer (e.g., Lightning)
- **Layer 2 approaches:**
  - **zkRollup, optimistic Rollup:** batch many Tx into a single Tx
- Sidechains: Polygon and others
- many other ideas ...

# (1) Payment channels (high level idea)

blockchain

\$100 held in channel (e.g., UTXO or DAPP)

Alice creates payment channel  
to Bob: value \$100

verify channel  
created correctly



Alice

**HTLC logic:** Hashed TimeLock Contract

Two ways to close channel:

- Tx with Alice sig: can close channel after 30 days, or
- Tx with Alice sig & Bob sig: close channel right away



Bob

# (1) Payment channels (high level idea)

blockchain

\$100 held in channel (e.g., UTXO)

Bob can sign Tx and close channel  
... but he would rather wait (up to 30 days)



Alice

pay Bob: 5\$

Tx: distribute funds: Alice: 95; Bob: 5

sig<sub>Alice</sub>



Bob

( off chain message! )

# (1) Payment channels (high level idea)

blockchain

\$100 held in channel (e.g., UTXO)



Alice

another payment: pay Bob: 15\$

Tx: distribute funds: Alice: 80; Bob: 20

sig<sub>Alice</sub>



Bob



# (1) Payment channels (high level idea)

blockchain

\$100 held in channel (e.g., UTXO)



Alice

another payment: pay Bob: 10\$

Tx: distribute funds: Alice: 70; Bob: 30

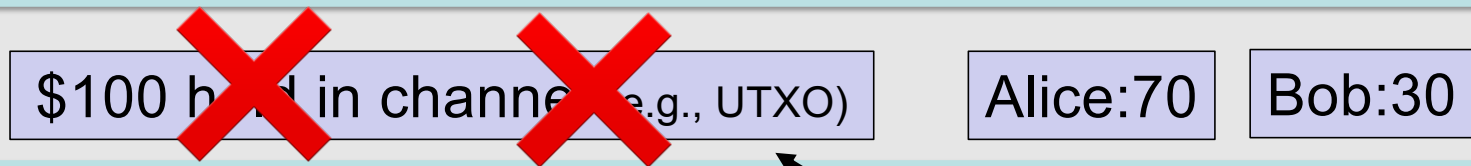
sig<sub>Alice</sub>



Bob

# (1) Payment channels (high level idea)

blockchain



either side can close channel  
(Alice only after 30 days)



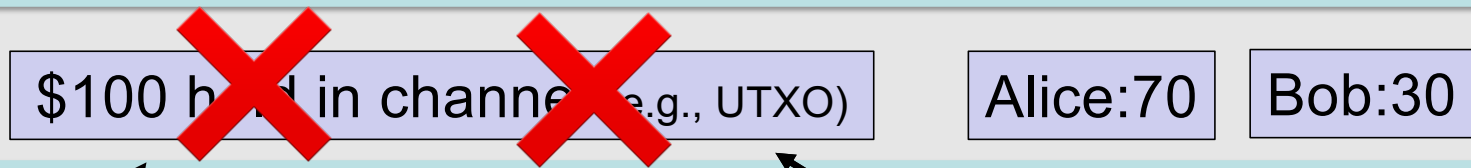
Alice



Bob

# (1) Payment channels (high level idea)

blockchain



either side can close channel  
(Alice only after 30 days)

Tx, sig<sub>A</sub>, sig<sub>B</sub>



Alice

main point: participants only touch chain  
when a channel is created or closed.

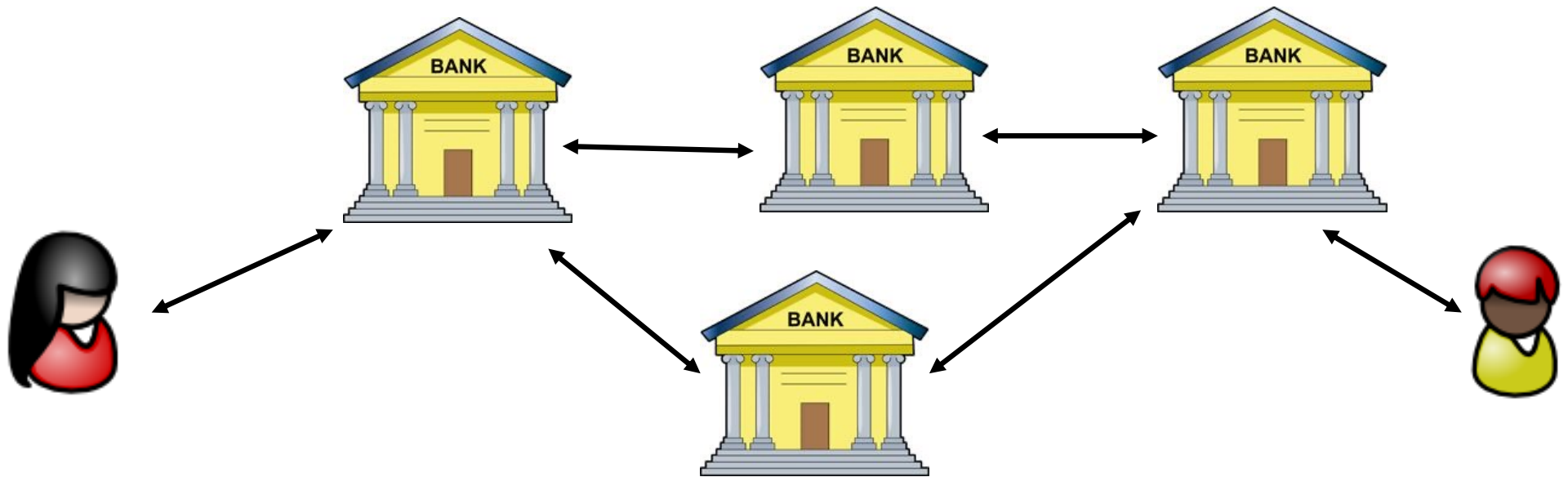


Bob

Bi-directional channels are also possible.

# Payment networks

Lots of bi-directional payment channels



Alice pays Bob by finding the cheapest route through the network  
⇒ while channels are open, nothing touches the blockchain

# The case of El Salvador

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June 24, 2021

8:18 PM PDT

Last Updated 2 months ago

**Technology**

## **Bitcoin to become legal tender in El Salvador on Sept 7, 2021**

3 minute read

Reuters

Payment channels are necessary to enable state-wide adoption

- Strike wallet: connects to the Bitcoin Lightning network

## (2) Scaling Ethereum Using Rollup

## (2) Scaling Ethereum Using Rollup

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Main tool: **SNARK** (much more on SNARKs later)

$C$ : a program that always terminates in  $\leq B$  steps

$x$ : public input to  $C$ ,  **$w$** : private input to  $C$

$(C, x, \mathbf{w})$



prover

short proof  $\pi$



$(C, x)$



verifier

## (2) Scaling Ethereum Using Rollup

### Main point:

Verifier's run time is  
\*much\* less than running  $C$

$x$ : public input to  $C$ ,

more on SNARKs

terminates in  $\leq B$  steps

$w$ : private input to  $C$

I am convinced  
prover knows  $w$   
s.t.  $C(x, w) = 1$

$(C, x, w)$



prover

short proof  $\pi$

$(C, x)$

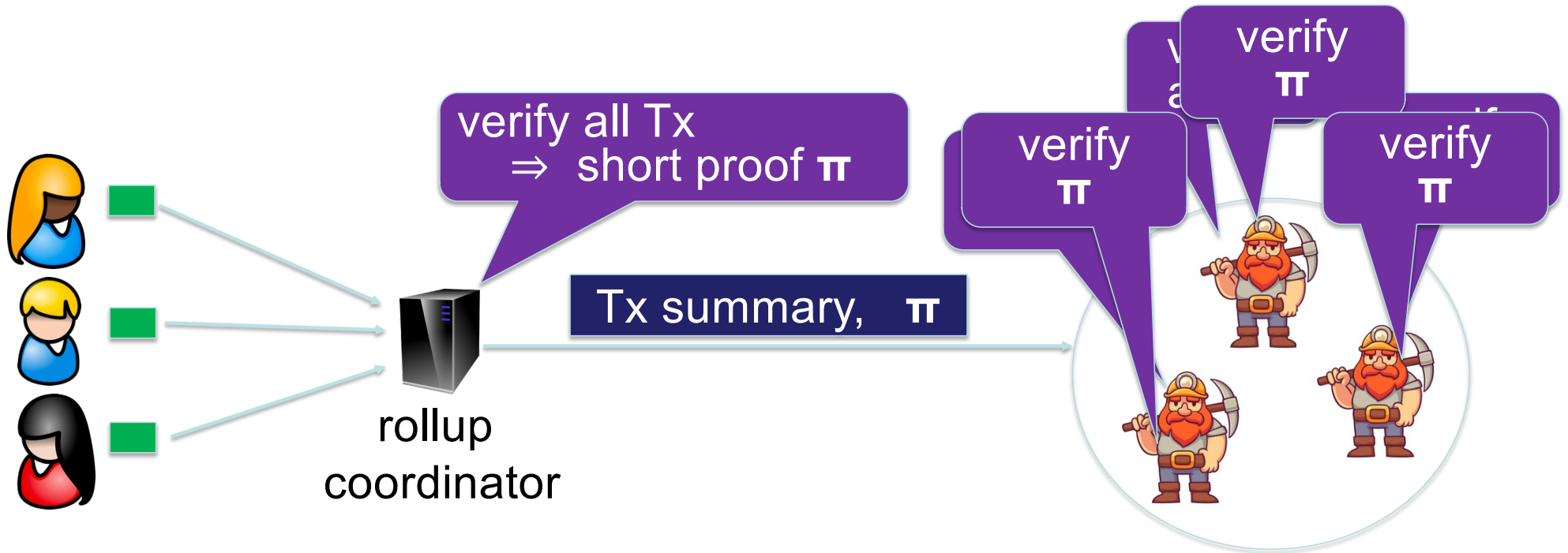


verifier



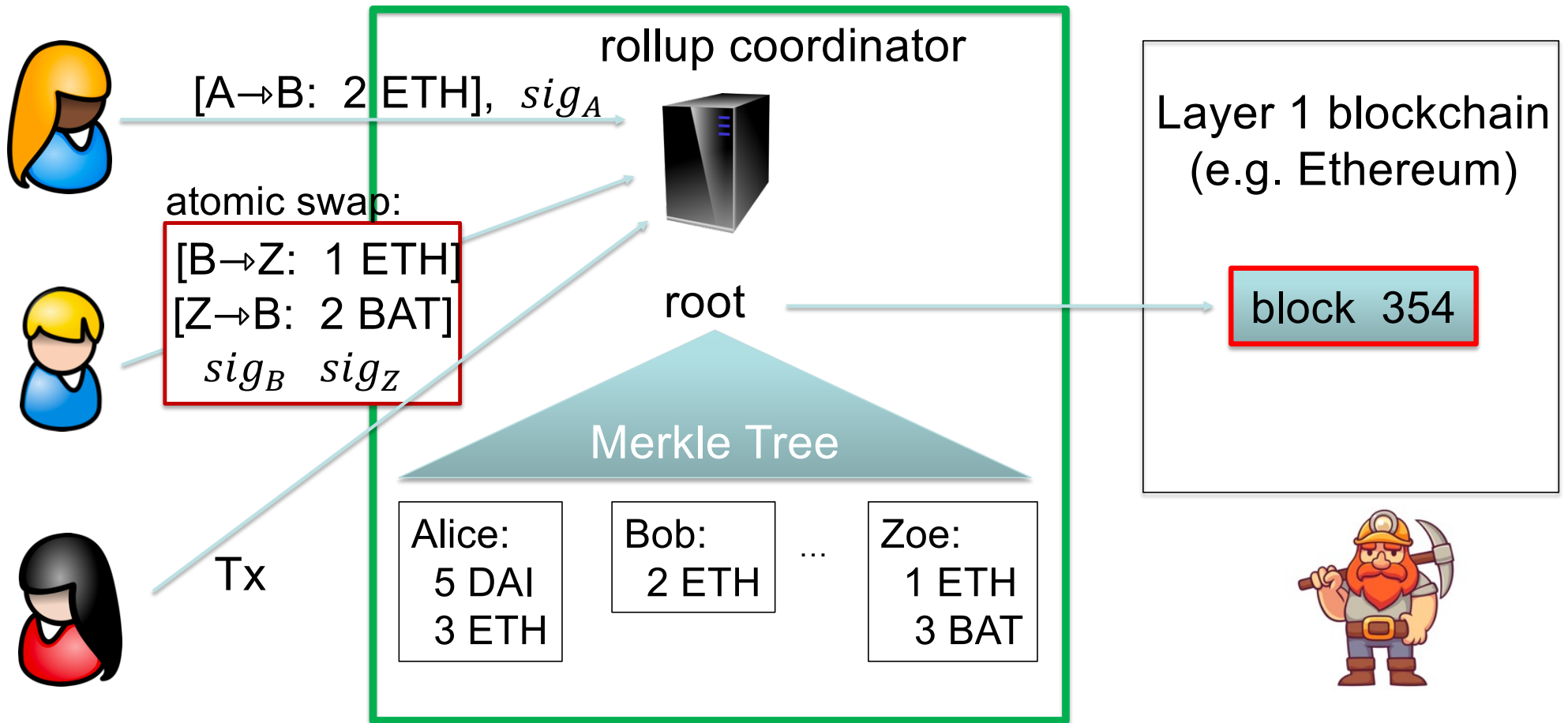
## (2) Rollup: zk and optimistic

Standard L1 chains: every miner must verify every posted Tx

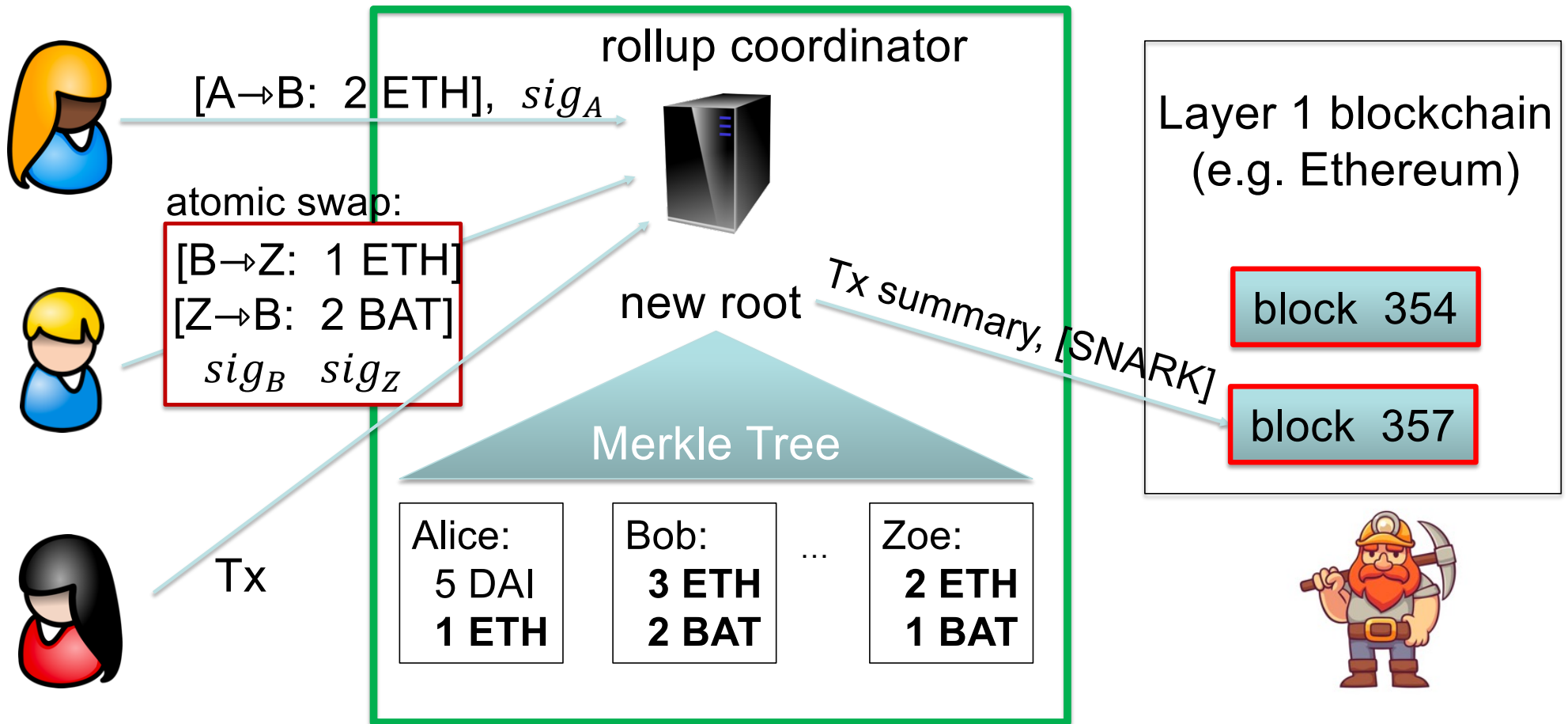


Rollup coordinator: compresses a thousand Tx into one on-chain proof (SNARK)

# zkRollup (simplified)



# zkRollup (simplified)



# Transferring assets to and from L2

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- Transactions within a Rollup system are easy:
  - Batch settlement on L1 network (e.g., Ethereum)
- Moving funds in to or out of Rollup system (L1  $\Leftrightarrow$  L2) is more expensive:
  - Requires posting more data on L1 network  $\Rightarrow$  higher Tx fees.
- Moving funds from one Rollup system to another (L2  $\Leftrightarrow$  L2)
  - Either via L1 network (expensive), or via a direct L2  $\Leftrightarrow$  L2 bridge (cheap)

# Migrating a project from L1 Ethereum to L2 zkRollup

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Upcoming development: **zkEVM** (e.g., MatterLabs and others).

## **Solidity compatibility:**

- Coordinator can produce a SNARK proof for the execution of a short Solidity program:
  - ⇒ easy to migrate a DAPP from L1 Ethereum to L2 zkRollup.
  - ⇒ reduced Tx fees and increased Tx rate compared to L1

# Optimistic Rollup (simplified) [e.g., Optimism, Arbitrum]

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Same principle as zkRollup, but no SNARK proof

Instead: coordinator posts Tx data on chain without a proof

- Then give a few days for validators to complain:

if a posted Tx is invalid  $\Rightarrow$

anyone can submit a **fraud proof** and win a reward,  
Rollup server gets slashed.

Benefit: simple full EVM compatibility, less work for server.

# Data availability: zkSync vs. zkPorter

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Is the coordinator a central point of failure? (centralization fears??)

Answer: No!

coordinator fails  $\Rightarrow$  users find another coordinator to produce proofs

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- Complication: new coordinator needs all current account information
  - How to get the data if the old coordinator is dead?
- Two solutions: zkSync and zkPorter. They work concurrently.

# Data availability: zkSync vs. zkPorter

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- **zkSync**: store all Tx summaries on the L1 blockchain (Ethereum)
  - L1 chain accepts Tx batch only if it includes summary of all Tx
  - Other coordinators can reconstruct L2 state from L1 blockchain
  - Downside: higher Ethereum Tx fees. Good for high value assets
- **zkPorter**: store Tx data on a new blockchain
  - maintained by a set of staked coordinators
  - Cheap off-chain storage, but lower guarantee than zkSync
- Customer can choose how coordinator will store its account.



That's it on this topic ...

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Next segment: interoperability

How to move assets from one chain to another



# Interchain Interoperability

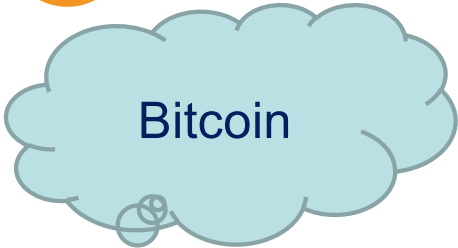
<https://defi-learning.org/>



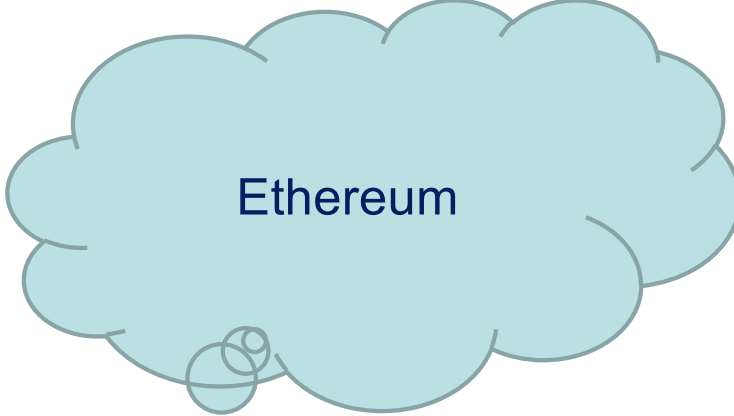
Solana



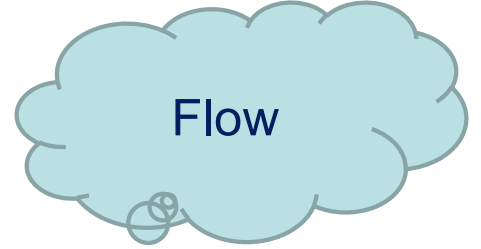
Serum  
DEX



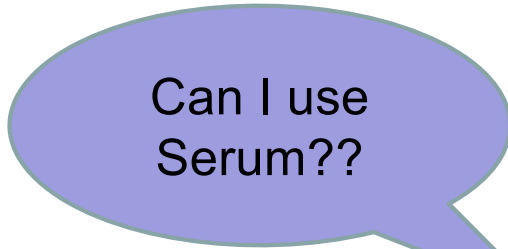
Bitcoin



Ethereum



Flow



Can I use  
Serum??



Polkadot



20 DOT

# Interoperability

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- **Interoperability:**

- a user owns funds or assets on one blockchain system.

Goal: enable the user to move funds and/or assets to another system.

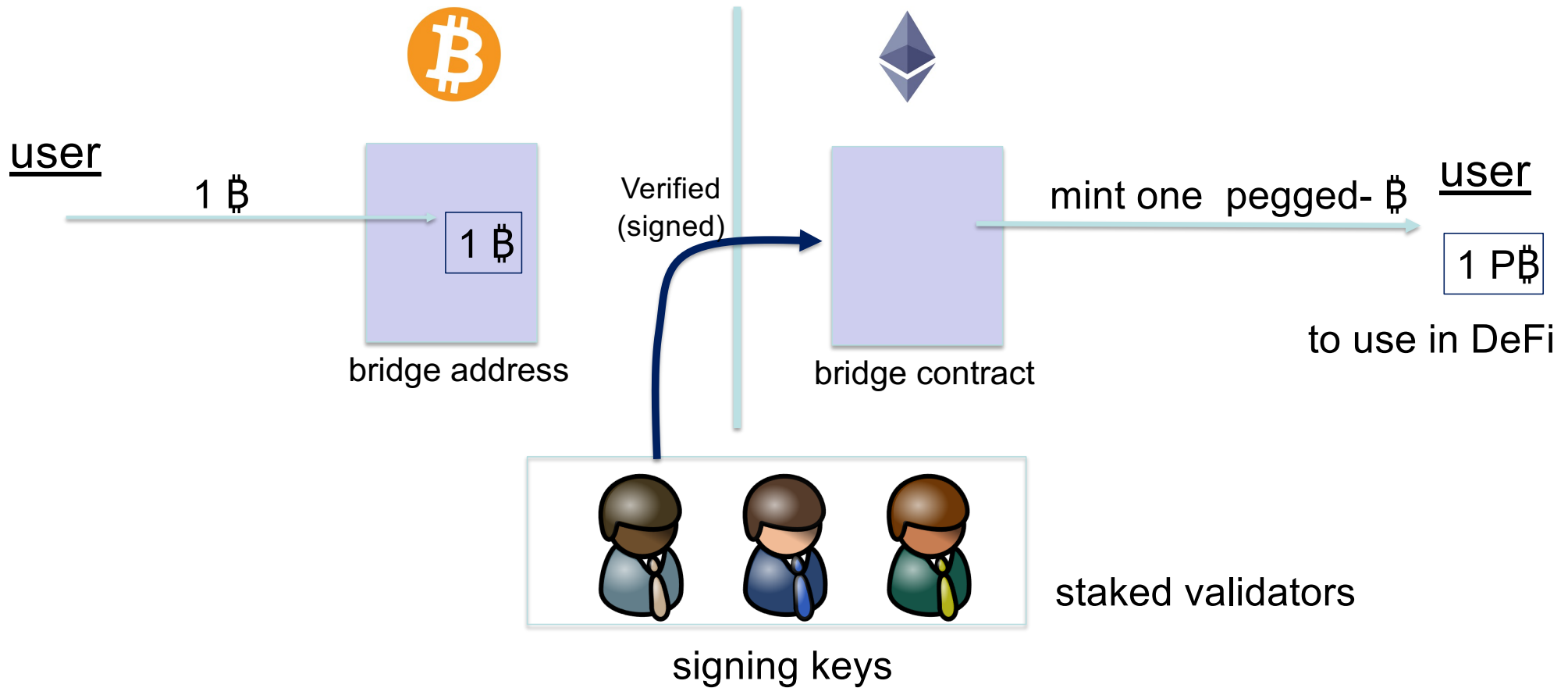
- **Composability:**

- enable a DAPP on one blockchain to call a DAPP on another

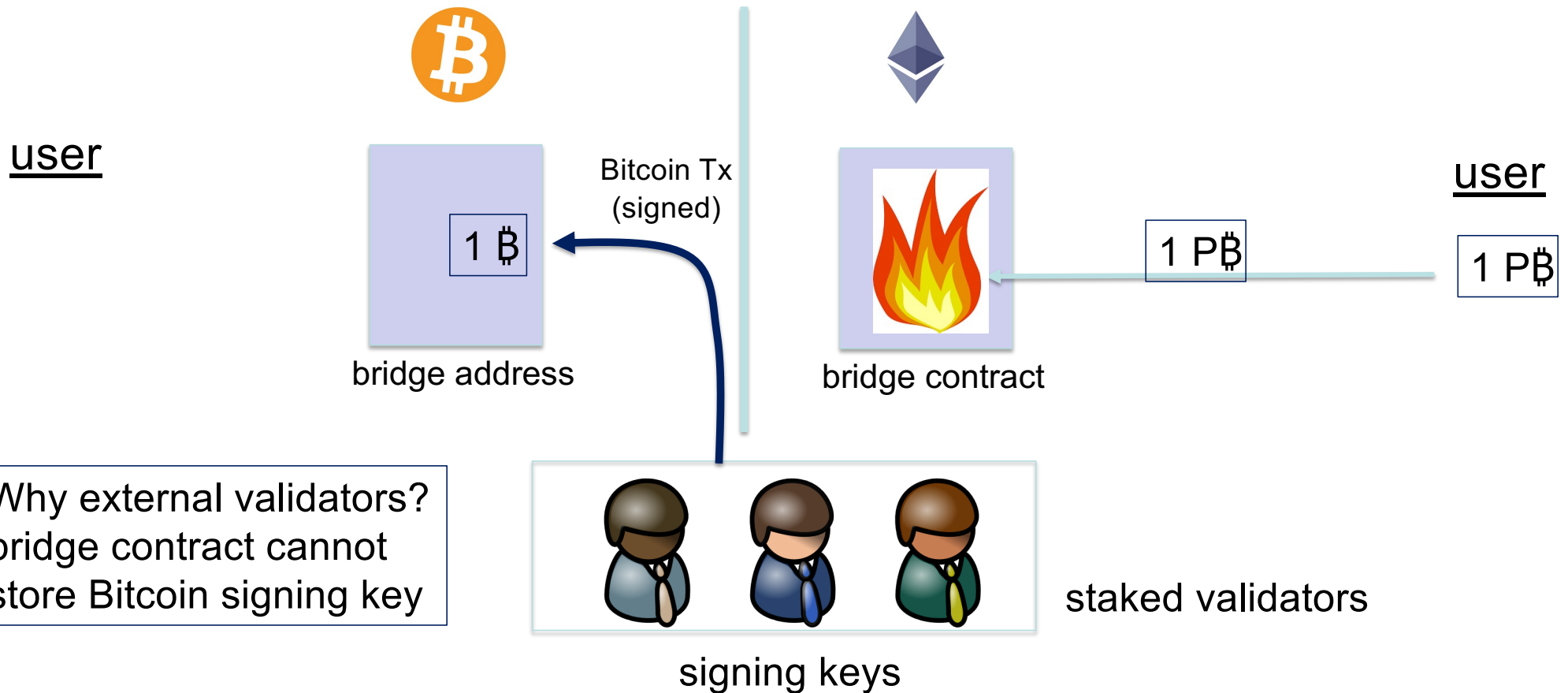
Both are easy if the entire world used Ethereum

- In reality: many blockchain systems that need to interoperate
- Several cross-chain protocols: XCMP, IBC, ...

# How to move assets? Building a federated bridge (simplified)



# How to move assets? Building a federated bridge (simplified)



# End of lecture: quick review

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## Cryptographic primitives:

- Hash functions: committing to large amounts of data
- Digital signatures: authorizing actions

## Scaling the blockchain

- Payment channels and Rollups

Interoperability: via bridges and pegged coins.



**END OF TOPIC**

<https://defi-learning.org/>